

Amendments to the Claims:

The following listing of claims will replace all prior versions, and listings, of claims in the application:

1. (Currently Amended) A method for generating a one-way function dependent on a one-way function H and a unique value d for a user, comprising:

holding in memory a function generation unique value s by a right issuer for the user;

creating a value generation unique value u in a unique value calculation unit from the function generation unique value s provided from the memory and the unique value d , the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to a token for the user;

creating by a hash value calculation unit a one-way function value $X(M)$ of a message M by applying the one-way function H to the value generation unique value u from the unique value calculation unit and the message M , where the one-way function value $X(M) = H(u_1 | M) | \dots | H(u_m | M)$;

holding a certificate C to prove a public key y paired with the one-way function value $X(M)$;

issuing a capability χ from the right issuer to the user, the capability χ representing a right of the user in association with the message M ; and

verifying the user from a the public key y and the capability χ by a right verifier.

2. (Original) The method for generating a one-way function according to claim 1, wherein the value generation unique value u is calculated by applying a one-way function G to the function generation unique value s and the unique value d .

3. (Original) The method for generating a one-way function according to claim 1, wherein the value generation unique value u is calculated by applying an encryption function E of a symmetric key to the function generation unique value s and the unique value d .

4. (Previously Presented) The method for generating a one-way function according to claim 1, wherein the one-way function value $X(M)$ of the message M is calculated by applying the one-way function H and an encryption function E of a symmetric key to the value generation unique value u and the message M .

5. (Currently Amended) A device for generating one-way function values that calculates a one-way function X dependent on a unique value d for a user, comprising:

means for inputting the unique value d ;

means for inputting a message M ;

means for holding a function generation unique value s by a right issuer for the user;

means for creating a value generation unique value u from the function generation unique value s from the holding means and the unique value d , the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to a token for the user;

means for creating a one-way function value $X(M)$ of the message M by applying a one-way function H to the value generation unique value u from the u -creating means and the message M , where the one-way function value $X(M) = H(u_1 | M) | \dots | H(u_m | M)$;

means for holding a certificate C to prove a public key y paired with the one-way function value $X(M)$;

means for issuing a capability χ from the right issuer to the user, the capability χ representing a right of the user in association with the message M; and

means for verifying the user from ~~a~~the public key y and the capability χ .

6. (Original) The device for generating one-way function values according to claim 5, wherein the process of calculating the value generation unique value u and the one-way function value $X(M)$ is difficult to observe from the outside.

7. (Currently Amended) A proving device for performing processing based on a private key for a user dependent on a message M, comprising:

means for inputting the message M, the message M including at least identifiers of private key processing algorithms;

means for holding a value generation unique value u for the user;

means for creating a one-way function value $X(M)$ of the message M by applying a one-way function H to the value generation unique value u from the holding means and the message M;

means for performing processing based on the private key and the one-way function value $X(M)$;

means for issuing a capability χ from the right issuer to the user, the capability χ representing a right of the user in association with the message M; and

means for verifying the user from a public key y and the capability χ ,

wherein the value generation unique value u is created from a function generation unique value s being held and provided by a right issuer and a unique value d for the user, the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to a token for the user, and the one-way function value $X(M) = H(u_1 \mid M) \mid \dots \mid H(u_m \mid M)$; and

wherein the identifiers in the message M enable the private key processing algorithms to be modified.

8. (Previously Presented) The proving device according to claim 7, wherein the calculation process in processing based on the value generation unique value u and the one-way function value $X(M)$ is difficult to observe from the outside.

9. (Original) The proving device according to claim 7, wherein the proving device is configured as a small portable operation device such as a smart card.

10. (Original) The proving device according to claim 7, wherein the proving device is configured as a module inside a CPU of the device.

11. (Previously Presented) The proving device according to claim 7, wherein the means for performing processing based on the private key comprises:

means for inputting a challenge c ;

means for calculating a response r from the challenge c and the one-way function value $X(M)$; and

means for outputting the response r .

12. (Previously Presented) The proving device according to claim 7, wherein the means for performing processing based on a private key comprises:

means for inputting a challenge c ;

means for generating a random number k ;

means for calculating a response r from the random number k , the challenge c , and the one-way function value $X(M)$; and

means for outputting the response r .

13. (Previously Presented) The proving device according to claim 7, wherein the means for performing processing based on a private key comprises:

means for generating a random number k ;

means for calculating a commitment w from the random number k ;

means for inputting a challenge c ;

means for calculating the response r from the random number k , the challenge c , and the one-way function value $X(M)$; and

means for outputting the response r .

14. (Previously Presented) The proving device according to claim 7, wherein the means for performing processing based on a private key comprises:

means for generating a random number k ;

means for calculating a commitment w from the random number k ;

means for outputting the commitment w ;

means for inputting a challenge c ;

means for calculating a response r from the random number k , the commitment w , the challenge c , and the one-way function value $X(M)$; and

means for outputting the response r .

15. (Original) The proving device according to claim 7, wherein the means for performing processing based on a private key performs multiplications and power operations of multiplicative groups on a finite field.

16. (Original) The proving device according to claim 7, wherein the means for performing processing based on a private key performs additions and scalar multiplication operations of elliptic curves on a finite field.

17. (Original) The proving device according to claim 7, wherein the means for performing processing based on a private key performs multiplicative residue operations and power residue operations modulo n , where n is a composite number that is difficult to factorize.

18. (Original) The proving device according to claim 7, wherein the message M includes use conditions and the means for inputting messages rejects message input if the use conditions included in the message M are not satisfied.

19. (Original) The proving device according to claim 7, wherein the message M includes private key processing parameters, and the means for performing processing based on a private key performs processing based on the private key processing parameters included in the message M.

20. (Currently Amended) A device for issuing a proving instrument T in accordance with a unique value d for a user, comprising:

means for inputting the unique value d;

means for holding a function generation unique value s by a right issuer for the user;

means for creating a value generation unique value u from the function generation unique value s from the holding means and the unique value d, the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to a token for the user;

means for writing the value generation unique value u from the u-creating means to the proving instrument T;

means for issuing the proving instrument T that includes a hash function X dependent on the unique value d;

means for issuing a capability χ from the right issuer to the user, the capability χ representing a right of the user in association with the message M; and

means for verifying the user from a public key y and the capability χ ,

wherein the proving instrument T holds the value generation unique value u, and upon input of a message M, creates a one-way function value $X(M)$ of the message M by

applying a one-way function H to the value generation unique value u and the message M to perform processing based on the one-way function value $X(M)$ expressed by $H(u_1 | M) | \dots | H(u_m | M)$.

21. (Previously Presented) An authentication method by which a right issuer issues rights to right recipients in association with a message M and a right verifier verifies the rights of the right recipients, the method comprising:

creating a value generation unique value u from a function generation unique value s being held and provided by a function generation unique value memory and a unique value d for a user corresponding to the right recipients, the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to a token for the user;

calculating a one-way function value $X(M)$ of the message M by a hash value generator by applying a one-way function H to the value generation unique value u and the message M , where the one-way function value $X(M) = H(u_1 | M) | \dots | H(u_m | M)$;

issuing a certificate C to prove a public key y paired with the one-way function value $X(M)$ to the right recipients by a certificate issuing unit;

presenting the certificate C from the right recipients to the right verifier;

performing processing by a private key processing unit based on the one-way function value $X(M)$;

verifying the certificate C by a certificate verification unit; and

verifying the processing by a private key processing verification unit based on the one-way function value $X(M)$ of the right recipients with a public key y proved by the certificate C .

22. (Original) The authentication method according to claim 21, wherein an identifier aid indicating an authentication type is included in the certificate C issued by the right issuer and the right verifier succeeds in verifying the certificate C only when the

authentication identifier aid included in the certificate C matches the type of authentication to be performed.

23. (Original) The authentication method according to claim 21, wherein use conditions are included in the certificate C issued by the right issuer and the right verifier succeeds in verifying the certificate C only when the use conditions included in the certificate C are satisfied.

24. (Previously Presented) A certificate issuing device for issuing a certificate C in accordance with a unique value d for a user and a message M, comprising:

means for inputting the unique value d;

means for inputting the message M;

means for holding a function generation unique value s by a right issuer for the user;

means for creating a value generation unique value u from the function generation unique value s from the holding means and the unique value d, the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to a token for the user;

means for creating a one-way function value $X(M)$ of the message M by applying a one-way function H to the value generation unique value u from the u-creating means and the message M, where the one-way function value $X(M) = H(u_1 \parallel M) \parallel \dots \parallel H(u_m \parallel M)$;

means for creating a public key y paired with the one-way function value $X(M)$;

means for issuing a certificate C to prove the public key y;

means for issuing a capability χ to the user, the capability χ representing a right of the user in association with the message M; and

means for verifying the user from the public key y and the capability χ .

25. (Previously Presented) An authentication device for performing authentication in accordance with a message M , comprising:

means for inputting the message M ;

means for holding a value generation unique value u for a user;

means for creating a one-way function value $X(M)$ of the message M by applying a one-way function H to the value generation unique value u from the holding means and the message M ;

means for performing processing based on the one-way function value $X(M)$;

means for holding a certificate C to prove a public key y paired with the one-way function value $X(M)$;

means for verifying the certificate C ;

means for issuing a capability χ to the user, the capability χ representing a right of the user in association with the message M ;

means for verifying the user from the public key y and the capability χ ; and

means for verifying processing based on a private key of the user with the public key y ,

wherein the value generation unique value u is created from a function generation unique value s being held and provided by a right issuer and the unique value d for the user, the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to a token for the user, and where the one-way function value $X(M) = H(u_1 \parallel M) \parallel \dots \parallel H(u_m \parallel M)$.

26. (Previously Presented) An authentication method by which a right issuer issues rights to right recipients in association with a message M and a right verifier verifies the rights of the right recipients, the method comprising:

creating a value generation unique value u from a function generation unique value s being held and provided by a function generation unique value memory and a unique value d for a user corresponding to the right recipients, the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to corresponding tokens for the right recipients;

calculating a one-way function value $X(M)$ of the message M by a hash value generator by applying a one-way function H to the value generation unique value u from the right issuer and the message M ;

issuing an access ticket t by an access ticket issuing unit determined from a private key x and the one-way function value $X(M)$ to the right recipients, where $X(M) = H(u_1 | M) | \dots | H(u_m | M)$;

performing processing by a private key processing unit based on the one-way function value $X(M)$;

converting the processing by a private key processing conversion unit based on the one-way function value $X(M)$ to the processing based on the private key x by the access ticket t ; and

verifying the processing by a private key processing verification unit based on the one-way function value $X(M)$ of the right recipients with a public key y paired with the private key x by the right verifier.

27. (Original) The authentication method according to claim 21, wherein an identifier aid indicating an authentication type is included in the message M .

28. (Previously Presented) An access ticket issuing device for issuing an access ticket in accordance with a unique value d for a user and a message M , comprising:

means for inputting the unique value d ;

means for inputting the message M ;

means for holding a function generation unique value s by a right issuer for the user;

means for creating a value generation unique value u from the function generation unique value s from the holding means and the unique value d , the value generation unique value u being provided as a series of m values where $u = (u_1, \dots, u_m)$ to a token for the user;

means for creating a one-way function value $X(M)$ of the message M by applying a one-way function H to the value generation unique value u and the message M , where the one-way function value $X(M) = H(u_1 | M) | \dots | H(u_m | M)$;

means for creating the access ticket t from a private key x and the one-way function value $X(M)$;

means for issuing the access ticket t ;

means for issuing a capability χ from the right issuer to the user, the capability χ representing a right of the user in association with the message M ; and

means for verifying the user from a public key y and the capability χ .

29. (Original) The access ticket issuing device according to claim 28, wherein the access ticket t is calculated as a difference $(x - X(M))$ between the private key x and the one-way function value $X(M)$.

30. (Original) The access ticket issuing device according to claim 28, wherein the access ticket t is calculated as a quotient $x/X(M)$ between the private key x and the one-way function value $X(M)$.

31. (Previously Presented) The access ticket generation device according to claim 28, wherein the unique value d for the user is (d_1, \dots, d_m) , the value generation unique value u is (u_1, \dots, u_m) and the one-way function value $X(M)$ is generated from bit concatenation $H(u_1 | M) | \dots | H(u_m | M)$ of the value of the one-way function H and has a desired bit length.

32. (Previously Presented) The access ticket generation device according to claim 31, wherein the value generation unique value (u_1, \dots, u_m) is found from $u_j = G(s_j \parallel d)$ obtained by applying a one-way function G to the function generation unique value $s = (s_1, \dots, s_m)$.

33. (Previously Presented) An authentication device for performing authentication for a user in accordance with a message M , comprising:

means for inputting the message M ;

means for holding a value generation unique value u for the user;

means for creating a one-way function value $X(M)$ of the message M by applying a one-way function H to the value generation unique value u from the holding means and the message M ;

means for performing processing based on the one-way function value $X(M)$;

means for holding an access ticket t determined from a private key x and the one-way function value $X(M)$;

means for converting the processing based on the one-way function value $X(M)$ to processing based on the private key x by the access ticket t ;

means for holding a public key y paired with the private key x ;

means for issuing a capability χ from the right issuer to the user, the capability χ representing a right of the user in association with the message M ;

means for verifying the user from the public key y and the capability χ ; and

means for verifying the processing based on the private key x with the public key y ,

wherein the value generation unique value u is created from a function generation unique value s being held and provided by a right issuer and a unique value d provided for the user, the value generation unique value u being provided as a series of m

values where $u = (u_1, \dots, u_m)$ to a token for the user, and where the one-way function value

$$X(M) = H(u_1 \parallel M) \parallel \dots \parallel H(u_m \parallel M).$$

34. (Original) The authentication device according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a challenge c with the access ticket t .

35. (Original) The authentication device according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a response r with the access ticket t .

36. (Original) The authentication device according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a response r with the access ticket t and a challenge c .

37. (Original) The authentication device according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a challenge c with a commitment w and means for updating a response r with the access ticket t and the challenge c .

38. (Original) The authentication device according to claim 33, wherein the means for converting the processing based on the private key comprises means for updating a challenge c with the access ticket t and a commitment w , and means for updating a response r with the commitment w .